

Scaling limits of k -ary growing trees

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Colloque JPS

The model

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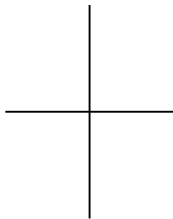
Fix an integer $k \geq 2$. We define a sequence $(T_n(k), n \geq 0)$ of random k -ary trees by the following recursion:

- $T_0(k)$ is the tree with a single edge and two vertices, a root and a leaf.
- given $T_n(k)$, to make $T_{n+1}(k)$, choose uniformly at random one of its edges, add a new vertex in the middle, thus splitting this edge in two, and then add $k - 1$ new edges starting from the new vertex.

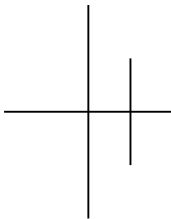
$\mathcal{T}_n(3)$ for $n = 0, 1, 2, 3$.



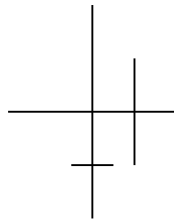
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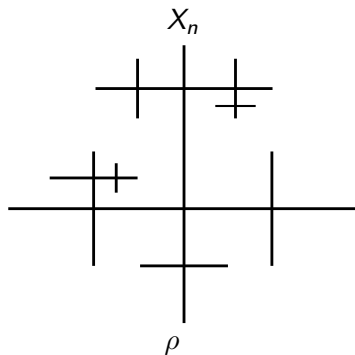


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$n = 10$



A few observations

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- $T_n(k)$ has n internal nodes, $kn + 1$ edges and $(k - 1)n + 1$ leaves
- When $k = 2$, we recover a well-known algorithm of Rémy, used to generate uniform binary trees. It is then well-known that the "size" of $T_n(2)$ is of order \sqrt{n} .

Size of $\mathcal{T}_n(k)$ and convergence

Distance between one point and the root

Let X_n be the point of $\mathcal{T}_n(k)$ corresponding to the leaf of $\mathcal{T}_0(k)$. Its distance to the root is of order $n^{1/k}$.

Lemma

$d(\rho, X_n)$ is of order $n^{1/k}$ as n goes to infinity, and in fact the quotient

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The proof of this lemma is simple. The process $(d(\rho, X_n), n \in \mathbb{Z}_+)$ is in fact a Markov chain: if we know that $d(\rho, X_n) = L$ then

$$d(\rho, X_{n+1}) = \begin{cases} L + 1 & \text{with probability } \frac{L}{(k-1)n+1} \\ L & \text{with probability } 1 - \frac{L}{(k-1)n+1} \end{cases}$$

Martingale methods then easily give us the existence of the limit we want.

The main convergence theorem

Let $\mu_n(k)$ be the uniform measure on the set of leaves of $T_n(k)$.

Theorem

As n tends to infinity, we have

$$\left(\frac{T_n(k)}{n^{1/k}}, \mu_n(k) \right) \xrightarrow{\mathbb{P}} (\mathcal{T}_k, \mu_k),$$

What kind of object is \mathcal{T}_k ?

\mathbb{R} -trees

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- for any continuous self-avoiding path $c: [0, 1] \rightarrow \mathcal{T}$, we have $c([0, 1]) = \varphi_{x,y}([0, d(x, y)])$, where $x = c(0)$ et $y = c(1)$.

Informally, an \mathbb{R} -tree is a connected union of line segments with no loops.

Graph-theoretical trees can be interpreted as \mathbb{R} -trees by considering each edge as a line segment and giving some length to each edge. In our case, the notation $\frac{T_n(k)}{n^{1/k}}$ means that we give to each edge of $T_n(k)$ a length equal to $n^{-1/k}$.

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However, the structure of an \mathbb{R} -tree can be very complex, notably because the leaves or the branch points can form a dense subset.

In practise, we will only want to look at rooted and measured trees: these are objects of the form $(\mathcal{T}, d, \rho, \mu)$ where ρ is a point on \mathcal{T} called the root and μ is a Borel probability measure on \mathcal{T} . Since d and ρ will never be ambiguous, we shorten the notation to (\mathcal{T}, μ) .

All our trees will also be compact.

The model
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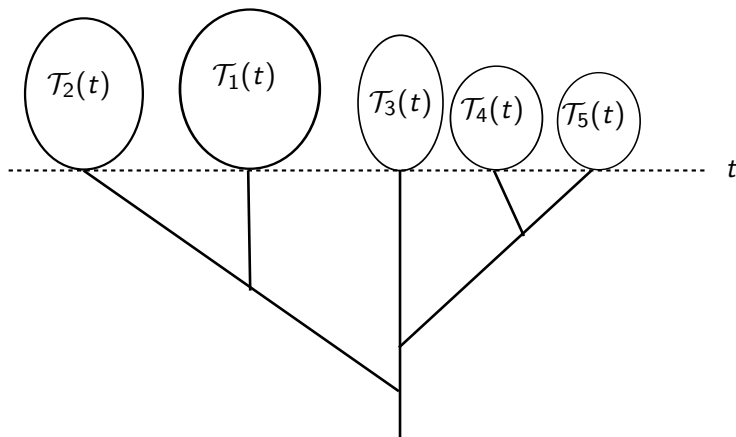
Stacking the trees
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Self-similar fragmentation trees

Self-similar fragmentation trees

Let $\alpha < 0$, and $(\mathcal{T}, d, \rho, \mu)$ be a compact random \mathbb{R} -tree.

For $t \geq 0$, we let $\mathcal{T}_1(t), \mathcal{T}_2(t), \dots$ be the connected components of $\{x \in \mathcal{T}, d(\rho, x) > t\}$.



Self-similar fragmentation trees

We say that (\mathcal{T}, μ) is a self-similar fragmentation tree with index α if, for all $t \geq 0$, conditionally on $\left(\mu(\mathcal{T}_i(s)); i \in \mathbb{N}, s \leq t\right)$:

- (Branching property) The subtrees $(\mathcal{T}_i(t), \mu_{\mathcal{T}_i(t)})$ are mutually independent.
- (Self-similarity) For any i , the tree $(\mathcal{T}_i(t), \mu_{\mathcal{T}_i(t)})$ has the same distribution as the original tree (\mathcal{T}, μ) , rescaled by $\mu(\mathcal{T}_i(t))^{-\alpha}$.

The notation $\mu_{\mathcal{T}_i(t)}$ means the measure μ conditioned to the subset $\mathcal{T}_i(t)$, which is a probability distribution.

Self-similar fragmentation trees

Linking these trees to the self-similar fragmentation processes of Bertoin shows that their distribution is characterized by three parameters:

- The index of self-similarity α .
- An *erosion coefficient* $c \geq 0$ which determines how μ is spread out on line segments.
- A *dislocation measure* ν , which is a σ -finite measure on the set

$$\mathcal{S}^\downarrow = \{\mathbf{s} = (s_i)_{i \in \mathbb{N}} : s_1 \geq s_2 \geq \dots \geq 0, \sum s_i \leq 1\}.$$

This measure determines how we allocate the mass when there is a branching point.

\mathcal{T}_k is a self-similar fragmentation tree

Theorem

The tree \mathcal{T}_k has the law of a self-similar fragmentation tree with:

- $\alpha = -\frac{1}{k}$.
- $c = 0$.
- *The measure ν_k is k -ary and conservative: it is supported on sequences such that $s_i = 0$ for $i \geq k + 1$ and $\sum_{i=1}^k s_i = 1$, and we have*

$$\nu(d\mathbf{s}) = \frac{(k-1)!}{k(\Gamma(\frac{1}{k}))^{(k-1)}} \prod_{i=1}^k s_i^{-(1-1/k)} \left(\sum_{i=1}^k \frac{1}{1-s_i} \right) \mathbf{1}_{\{s_1 \geq s_2 \geq \dots \geq s_k\}} d\mathbf{s}$$

Fractal dimension of \mathcal{T}_k

Corollary

The Hausdorff dimension of \mathcal{T}_k is almost surely equal to k .

This is a consequence of well-known results on fragmentation trees.

Stacking the trees

The model
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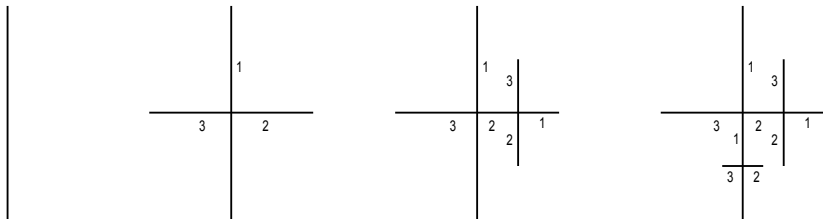
Stacking the trees
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Labelling the edges

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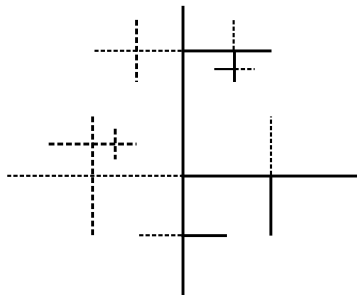
Each step of the algorithm creates k new edges. We give them labels 1 to k the following way:

- The upper half of the edge which was split in two is labeled 1
- The other new edges are labeled $2, \dots, k$.



$T_n(k')$ inside $T_n(k)$

Consider an integer $k' < k$. Let $T_n(k, k')$ be the subset of $T_n(k)$ where we have erased all edges with labels $k' + 1, k' + 2, \dots, k$ and all their descendants.



If we call I_n the number of internal nodes which are in $T_n(k, k')$, then one can check that:

- Conditionally on I_n , $T_n(k, k')$ is distributed as $T_{I_n}(k')$.

$\mathcal{T}_{k'}$ inside \mathcal{T}_k

One can show that the sequence $\frac{I_n}{n^{k'/k}}$ converges a.s. to a random variable M . As a consequence we obtain:

Proposition

$$\frac{T_n(k, k')}{n^{1/k}} \xrightarrow{\mathbb{P}} M \mathcal{T}_{k, k'}$$

where $\mathcal{T}_{k, k'}$ is a version of $\mathcal{T}_{k'}$ hidden in \mathcal{T}_k , and is independent of M .

More on the stacking

- It is in fact possible to extract directly from \mathcal{T}_k a subtree distributed as $\mathcal{T}_{k'}$, without going back to the finite case: at every branch point of \mathcal{T}_k , select only k' of the k branches at random with a well-chosen distribution.

Thank you!